

# Solving lower/upper bound approaches of limit analysis by an interior-point method for convex programming

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- ▶ demonstrate that large problems can be solved by a matlab-based implementation
- ▶ show the first result of a domain decomposition-like technique

## Statement of the static Limit Analysis problem

An infinite bar is compressed under two rough rigid plates. [▶ Figure](#)

A quarter of the section is meshed in triangular Lagrange p1 elements, with appropriate symmetry and boundary conditions. Solving the static problem leads to maximize a linear function of the stresses variables  $(\sigma_x, \sigma_y, \sigma_{xy})$  of each triangle's apex, under linear equalities constraints (equilibrium, continuity, symmetry and boundary conditions), and one non-linear inequality per apex. The latter depends on the selected criterion. For example, with the Mises criterion:

$$(\sigma_x - \sigma_y)^2 + (2\sigma_{xy})^2 \leq (2c)^2.$$

[▶ cont.](#)[▶▶ tests](#)

The solution is a lower bound for the Limit Analysis problem

- Introduction

- Statement of the problem: figure

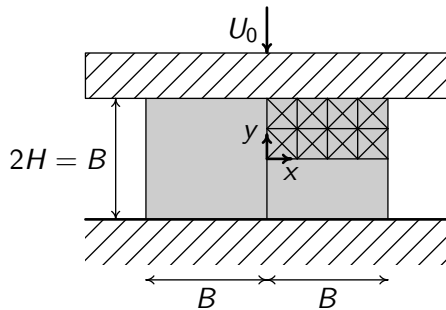


Figure: Compression of a bar between rough rigid plates





## Typical optimization problems from static Limit Analysis

General form of the static Limit Analysis mechanical optimization problems :

$$\begin{aligned} \max \quad & c^T x \\ \text{s.t.} \quad & Ax = b, \\ & g(x) \leq 0, \end{aligned}$$

where

- ▶  $c, x \in \mathbb{R}^n$ ,  $b \in \mathbb{R}^m$ ,  $A \in \mathbb{R}^{m \times n}$
- ▶  $g = (g_1, \dots, g_p)$  is a vector-valued function of  $p$  convex numeric functions  $g_i$ .

This problem is convex, both equality and inequality constrained, potentially from medium to large scale, sparse.

IP methods are particularly well suited for this kind of problem

- └ A solution approach for the convex programming problem

- └ The barrier problem

## Transformation in a barrier problem

We have adapted an IP algorithm proposed originally by VIAL (1992) for convex programming problems

The algorithm is of the type “primal-dual interior point method”.

The development of the algorithm is as follows:

The original problem, is transformed in an unconstrained “barrier problem”, with a parameter  $\mu > 0$ , the “barrier parameter”:

$$\begin{aligned} \max \quad & c^T x + \mu \sum_{i=1}^p \ln(s_i) \\ \text{s.t.} \quad & Ax = b, \\ & g(x) + s = 0. \end{aligned}$$

└ A solution approach for the convex programming problem

└ The KKT system

## The optimality condition for the barrier problem

The KKT conditions are:

$$\begin{aligned} -c + A^T w + \left( \frac{\partial g}{\partial x} \right)^T y &= 0 = F_d(x, y, w, s), \\ Ax - b &= 0 = F_{p_1}(x, w, y, s), \\ g(x) + s &= 0 = F_{p_2}(x, w, y, s), \\ YSe - \mu e &= 0 = F_c(x, w, y, s), \end{aligned}$$

where  $w \in \mathbb{R}^m$ ,  $y \in \mathbb{R}^p$  and  $Y, S$  are the diagonal matrices associated to  $y$  and  $s$  respectively.  $e \in \mathbb{R}^p$  is a vector of ones.

- └ A solution approach for the convex programming problem

- └ The Newton system

## The Newton system

The KKT conditions for the barrier problem is expressed as:

$F = (F_d, F_{p_1}, F_{p_2}, F_c) = 0$  (the Newton system).

The following linear system must be solved:

$$\begin{bmatrix} H_0 & A^T & \left(\frac{\partial g}{\partial x}\right)^T & 0 \\ A & 0 & 0 & 0 \\ \frac{\partial g}{\partial x} & 0 & 0 & I \\ 0 & 0 & S & Y \end{bmatrix} \begin{bmatrix} dx \\ dw \\ dy \\ ds \end{bmatrix} = \begin{bmatrix} -F_d \\ -F_{p_1} \\ -F_{p_2} \\ -F_c \end{bmatrix}.$$

Given that  $g$  is convex,  $H_0 = \sum_{i=1}^p y_i \frac{\partial^2 g_i}{\partial x^2}$  is positive semi-definite, and in some cases positive definite.

- └ A solution approach for the convex programming problem

- └ The equilibrium system

## Solving the Newton system: the equilibrium system

Some row and column reordering and a “block-elimination” reduction lead to:

$$\begin{bmatrix} Y & S & 0 & 0 \\ 0 & -Y^{-1}S & \frac{\partial g}{\partial x} & 0 \\ 0 & 0 & H_0 + \left(\frac{\partial g}{\partial x}\right)^T YS^{-1} \frac{\partial g}{\partial x} & A^T \\ 0 & 0 & A & 0 \end{bmatrix} \begin{bmatrix} ds \\ dy \\ dx \\ dw \end{bmatrix} = \begin{bmatrix} -F_c \\ -F_{p_2} + Y^{-1}F_c \\ -F_d - \left(\frac{\partial g}{\partial x}\right)^T YS^{-1}r \\ -F_{p_1} \end{bmatrix},$$

with  $r = -F_{p_2} + Y^{-1}F_c$ .

└ A solution approach for the convex programming problem

└ The equilibrium system

## Solving the Newton system: the equilibrium system (cont.)

Let us define  $H = H_0 + \left(\frac{\partial g}{\partial x}\right)^T YS^{-1} \frac{\partial g}{\partial x}$ . Thus we first have to solve the following system:

$$\begin{bmatrix} H & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} dx \\ dw \end{bmatrix} = \begin{bmatrix} -F_d - \left(\frac{\partial g}{\partial x}\right)^T YS^{-1} r \\ -F_{p_1} \end{bmatrix}.$$

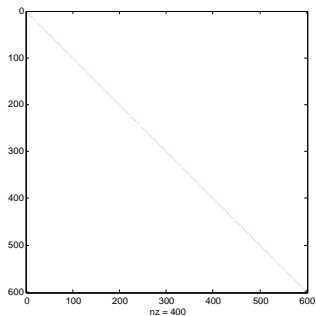
- ▶ This kind of system is known as an “equilibrium system”.
- ▶ The equilibrium system is symmetric, never definite...

NB: To achieve the decrease of  $\mu$  and computing the search direction  $dz$  along  $\mathcal{C}$ , we have implemented the Mehrotra predictor-corrector algorithm.

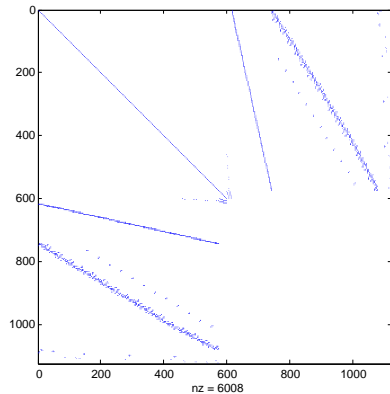
- └ A solution approach for the convex programming problem
  - └ The equilibrium system

## Solving the equilibrium system

The matrix  $H$



The matrix of the equilibrium system

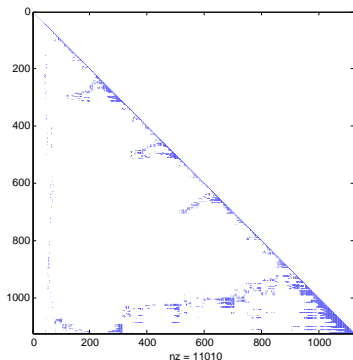


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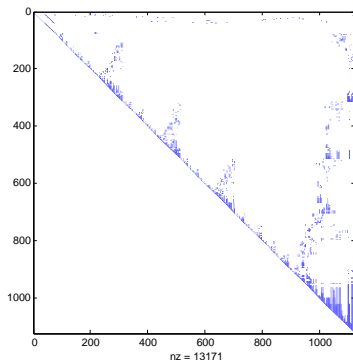
## Solving the equilibrium system

Use of a specific method, or LU factorization

The factor  $L$  of LU decomposition



The factor  $U$  of LU decomposition





## Application to the Mises and Gurson Criteria

We already gave an example of a plasticity criterion : the Mises one. ◀

Another criterion: Gurson. The exact solution is not known in this case. Hereafter,  $f = 0.16$ .

$$(\sigma_x - \sigma_y)^2 + (2\sigma_{xy})^2 + 8c^2 f \cosh \frac{(\sigma_x + \sigma_y)}{2k} \leq 4c^2(1 + f^2).$$

- ▶ It gives rise to another convex programming, **not** a conic programming problem.
- ▶ A series of tests on the preceding mechanical system, involving Mises and Gurson criteria, were performed.

- └ Computational experiments
  - └ Test problems

## Computational Statistics

Experiments with Matlab 6.5.1.

On Apple Macintosh dual G5 2.5Ghz, 4.5GB of ram, under MacOSX. Only one processor used, and 2GB of ram (32bits code).

$N_{tr}$	Vars.	Constraints		Mises			Gurson		
		Lin.	Conv.	Res.	Iter.	Time	Res.	Iter.	Time
800	7 440	6 340	2 480	2.41346	18	70s	1.64768	14	18s
7 200	65 520	56 220	21 840	2.42270	18	12m 21s	1.64950	19	44m
20 000	181 200	155 700	60 400	2.42465	20	1h 32 m	1.64989	27	7h 24m

Table: Mises and Gurson criteria : comparison.

NB: the exact solution of this problem is known for Mises: 2.42768.

└ Computational experiments

└ Static dual method (linear continuous velocity fields)

## The kinematic problem

- ▶ The kinematic problem aims at producing an upper bound for the Limit Analysis problem.
- ▶ Difficulty: following the usual way to solve this problem, one has to integrate the dissipated power  $\pi$ , which is sometimes very complicated to compute, if not possible
- ▶ Hence the interest of a kinematic method which does not use the dissipated power expression, only the plasticity criterion expression.

└ Computational experiments

└ Static dual method (linear continuous velocity fields)

## A new kinematic approach

The external power can be expressed as  $Q \cdot q$ , where :

- ▶  $Q$  the load vector,
- ▶  $q = q(u)$  the generalized velocity vector, with  $u$  kinematically admissible (KA).

### Virtual Power Principle

$Q$  and  $\sigma^*$  are in equilibrium if, for all KA vectors  $u$  :

$$Q \cdot q = \int_V \sigma : v \, dV$$

where  $v$  is the strain rate tensor (which depends on  $u$ ).

- Computational experiments

- Static dual method (linear continuous velocity fields)

## Finite element discretization

Assuming a 1-dimensional loading problem  $Q^*$ , for sake of simplicity, and the velocities  $u$  **linear** and **continuous**:

$$\left. \begin{aligned} qQ^* &= \{u\}^T \{\beta\} Q^* \\ \int_V \sigma^* : v \, dV &= \{u\}^T [\alpha] \{\sigma^*\} \end{aligned} \right\} \Rightarrow \{u\}^T ([\alpha] \{\sigma^*\} - \{\beta\} Q^*) = 0 \quad \forall \{u\} \text{ KA}$$

It leads to the following problem, in the case of the compressed bar with  $q = U_0$ :

$$\begin{aligned} \max \quad & qQ^* \\ \text{s.t.} \quad & [\alpha] \{\sigma^*\} - \{\beta\} Q^* = 0, \\ & f(\sigma^*) \leq 0, \quad \sigma^* \text{ constant in each finite element} \\ & + \text{ limit, symmetric and loading linear conditions.} \end{aligned}$$

└ Computational experiments

└ Static dual method (linear continuous velocity fields)

A KKT condition on the discretized problem:

$$-c + A^T w + \left( \frac{\partial f}{\partial x} \right)^T y = 0,$$

where :

$$A = \left[ [\alpha], -\{\beta\} \right]^T, \quad x = \left\{ \left\{ \sigma^* \right\} \right\} \left\{ Q^* \right\}.$$

- ▶  $w$ : dual variables associated to linear constraints;
- ▶  $y$ : dual variables associated to non linear constraints;
- ▶  $c$ : coefficients of the objective function.

- └ Computational experiments

- └ Static dual method (linear continuous velocity fields)

By analyzing the structure of the following equations:

$$\begin{aligned}\{u\}^T \left[ [\alpha] \{\sigma^*\} - \{\beta\} Q^* \right] &= 0, \\ -\{c\}^T + \{w\}^T A + \{y\}^T \left\{ \frac{\partial f}{\partial \sigma} \right\} &= 0,\end{aligned}$$

it can be proved that the field  $u = -w$  is admissible (ie KA and PA). Hence the method is rigorously kinematic, requiring *only* the plasticity criterion as information about the material.

└ Computational experiments

└ Static dual method (linear continuous velocity fields)

## Application: compressed bar, plane strain

- ▶ for this kind of problem (Gurson), it has been possible to use a Cholesky factorisation to solve the main system
- ▶ For the largest problems, the main linear system had to be slightly perturbed (diagonal perturbation of  $10^{-8}$  on  $H$ )

▶ Figure



- Computational experiments

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▶ Figure

$N_{tr}$	Variables	Constraints		Lin. cont. u	
		Lin.	Conv.	Opt. Value	Time.
400	2 403	790	801	1.6779	4s
3 200	9603	3 180	3 201	1.6655	44s
7 200	7 201	21 603	7 170	1.6611	2m 19s
20 000	60 003	19 950	20 001	1.6572	16m 3s

**Table:** The compressed bar and the *Gurson* material - *kinematic* results for  $f = 0.16$  using *linear continuous* velocity fields

- └ Computational experiments

- └ Static dual method (linear continuous velocity fields)

Static bound for a Gurson material ( $f = 0.16$ ): 1.6499

$$1.6499 < \frac{F}{2Bk} < 1.6572$$

Only 0,7% of gap between the two bounds.

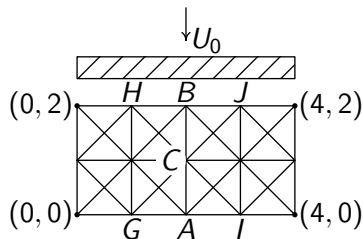
## For bigger problems: a “divide and conquer” strategy

- ▶ When problems get too large, “out of memory” occurs within Matlab and matrices are increasingly bad-conditioned.
- ▶ Hence the idea: splitting the problem in two (or more).
- ▶ It happens to be possible in the static-dual algorithm, because of the mechanical meaning of all numerical variables in this problem.

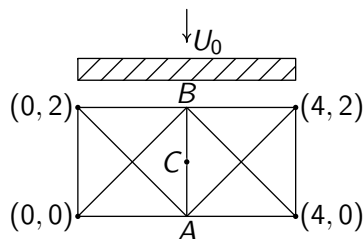
## The general idea on a simple example

Purpose : solving a static-dual kinematic problem with linear continuous velocity fields.

- ▶ The loading vector:  $Q = F$
- ▶ The generalized velocity vector:  $U_0$



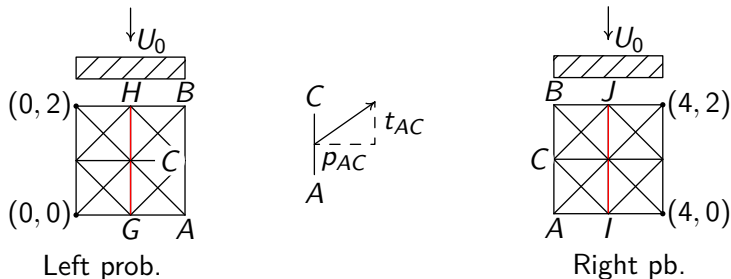
$4 \times 2$  problem (to solve)



Starting  $2 \times 1$  problem

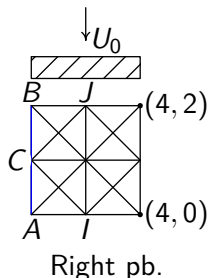
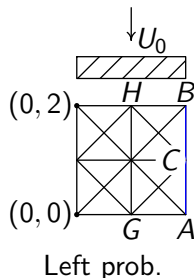
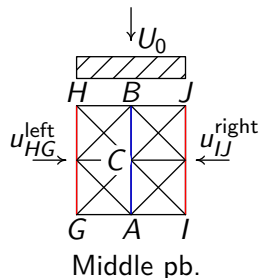
$$u_C = (u_A + u_B) / 2$$

## First iteration



- ▶ The loading vector:  $Q = (F, p_{AC}, t_{AC}, p_{CB}, t_{CB})$
- ▶ The velocity vector:  $q = (U_0, (u_A + u_B)^T/2, (u_B + u_C)^T/2)$ , where vectors  $u_A, u_B, u_C$  are collected on the starting problem
- ▶ The sum of the objective values of these problems is a kinematic bound

## Subsequent iterations



- ▶ The sum of the objective values of the latter two problems is another kinematic bound, noticeably lower than the previous one. This bound improves steadily if we iterate this process.
- ▶ At each iteration, only the coefficients of the objective function change.

## First experiments

- ▶ On the compressed bar
- ▶ Performed on a PowerbookG4, 1.33Ghz, 2Go of RAM.

	Original problem		Splitted problem		
Size	Result	Time	Global Iter.	Result	Time
$16 \times 8$	2.53284	n.s.	5	2.53303	n.s.
$32 \times 16$	2.48753	59s	2	2.48597	68s
$64 \times 32$	2.45833	1030s	2	2.45957	720s

- ▶ The results are more and more accurate as iterations go on.
- ▶ However, only a few iteration are necessary, in a reasonable amount of time.

## Conclusion and future work

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- ▶ LA problems are also interesting as test-bed problems for people working on large sparse symmetric structured systems of linear equations, positive definite systems and indefinite ones.

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- ▶ LA problems are also interesting as test-bed problems for people working on large sparse symmetric structured systems of linear equations, positive definite systems and indefinite ones.
- ▶ The use of domain decomposition-like techniques makes parallel processing attractive.