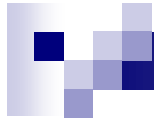




Frequent Pattern Mining

Toon Calders

University of Antwerp



Summary

- Frequent Itemset Mining
- Algorithms
- Constraint Based Mining
- Condensed Representations

Frequent Itemset Mining

■ Market-Basket Analysis

transaction identifier

TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

items

transaction

Frequent Itemset Mining

- support(I): number of transactions “containing I”

TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

Support(BC) = 3

Support(ACD) = 2

Frequent Itemset Mining Problem

Given D, minsup

Find all sets I with $\text{support}(I) \geq \text{minsup}$

TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

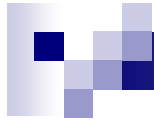
minsup=2

{}, A, B, C, D,
AC, AD, BC, BD, CD,
ACD

Why?

- Important component in mining algorithms
- Sufficient statistics for interestingness measures
 - Confidence $X \rightarrow Y$: $\text{Support}(XY) / \text{Support}(X)$
 - Contingency tables (correlation, χ^2)

	X	$\neg X$
Y	$s(XY)$	$s(Y) - s(XY)$
$\neg Y$	$s(X) - s(XY)$	$s(\{\}) - s(X) - s(Y) + s(XY)$



Summary

- Frequent Itemset Mining
- Algorithms
- Constraint Based Mining
- Condensed Representations



Algorithms

- There exist hundreds of algorithms that solve FIM (or related problems)
 - AIS, Apriori, AprioriTID, AprioriHybrid, FPGrowth, FPGrowth*, Eclat, dEclat, Pincersearch, ABS, DCI, kDCI, LCM, AIM, PIE, ARMOR, AFOPT, COFI, Patricia, MAXMINER, MAFIA, NDI-ALL, ...



Algorithms

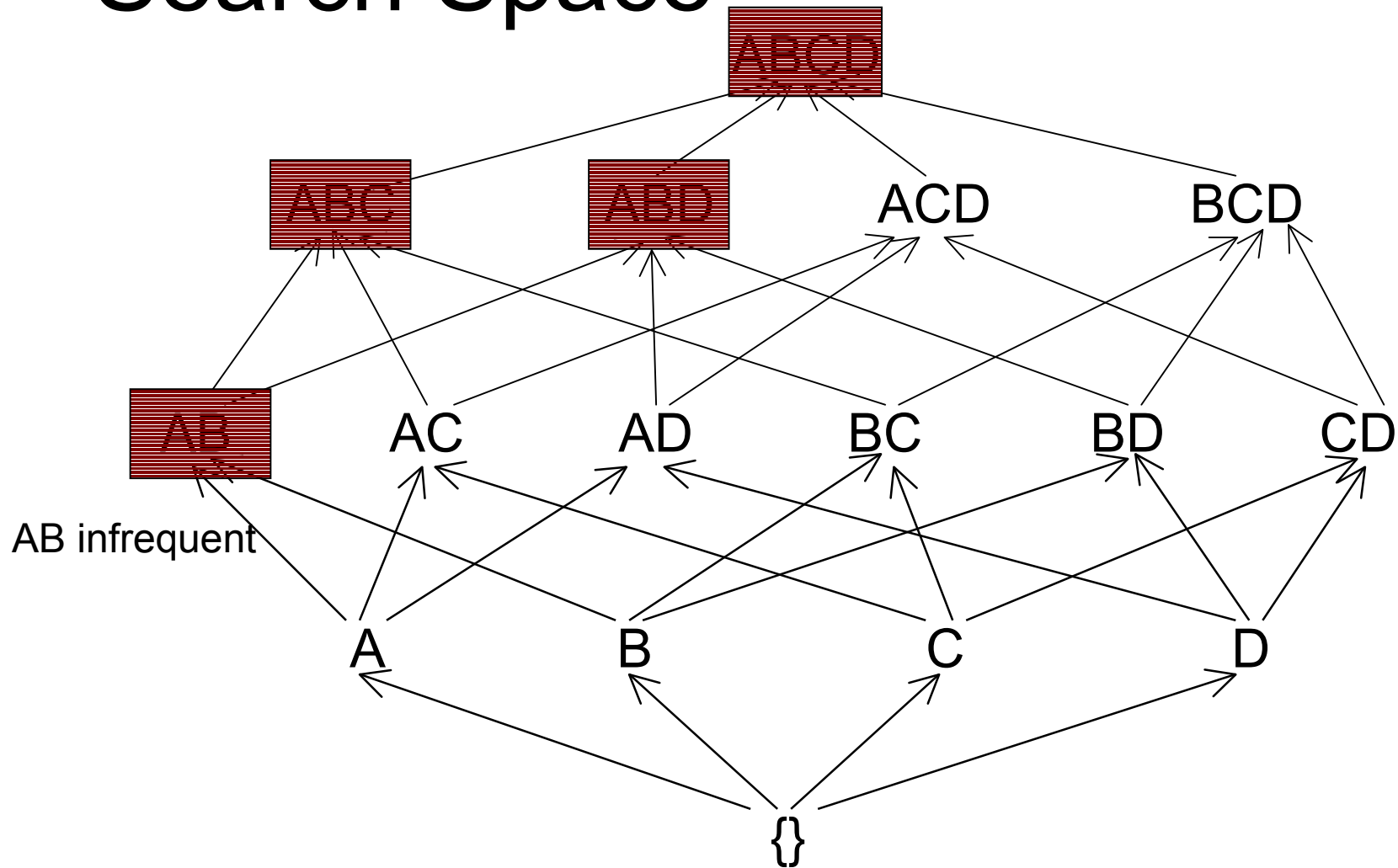
- There exist hundreds of algorithms that solve FIM (or related problems)
 - Concentrate on the most important pruning principle:
 - Monotonicity
- and the two main search strategies:
- Breadth-first
 - Depth-first



Monotonicity Principle

- If $I \subseteq J$, then $\text{support}(I) \geq \text{support}(J)$
- Therefore, if I is infrequent, then all its supersets are infrequent as well.
- All FIM algorithms rely heavily on this principle to prune large parts of the search space.

Search Space





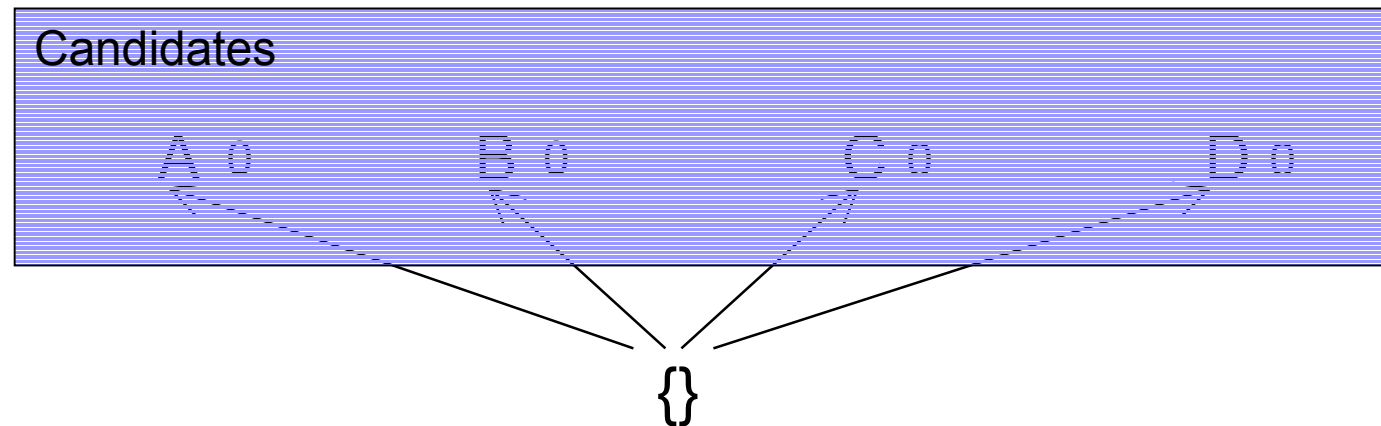
Levelwise Algorithm

- Exploits monotonicity as much as possible.
- Search Space is traversed bottom-up, level by level
- Support of an itemset is only counted in the database if all its subsets were frequent.

TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

Apriori

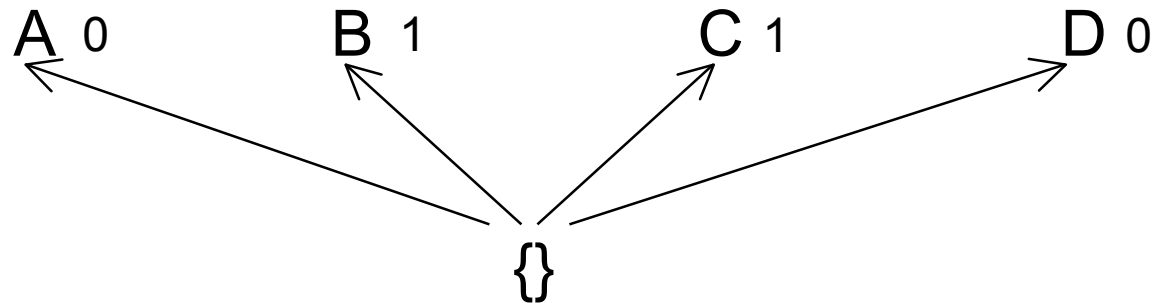
minsup=2



TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

Apriori

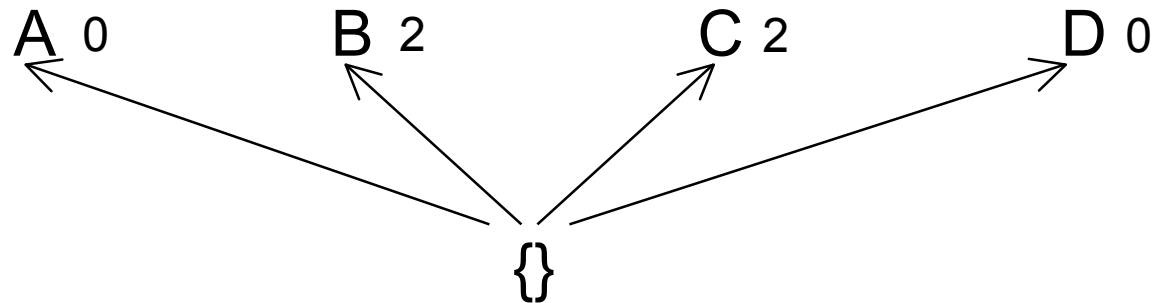
minsup=2



TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

Apriori

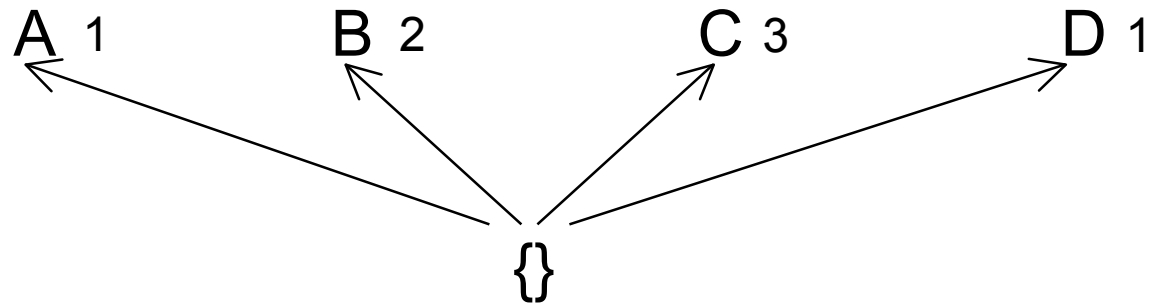
minsup=2



TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

Apriori

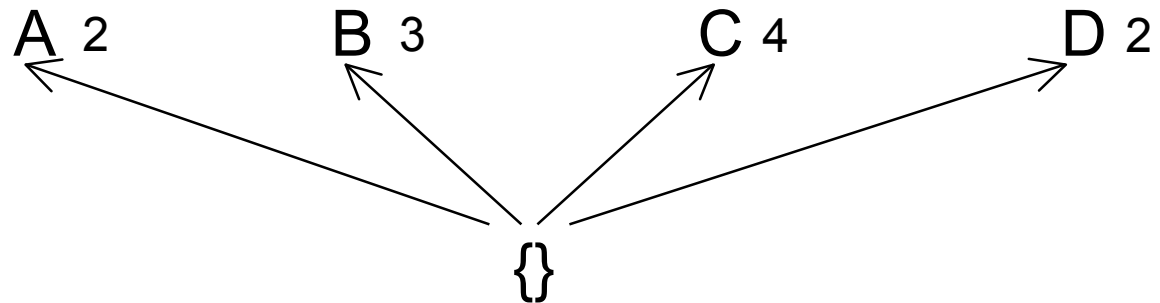
minsup=2



TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

Apriori

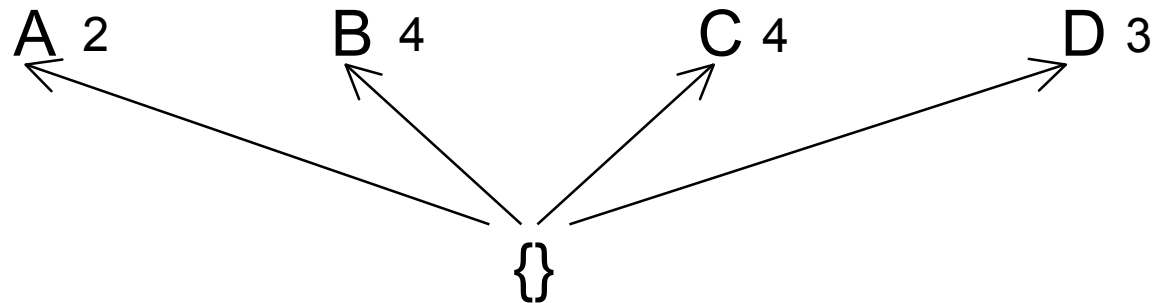
minsup=2



TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

Apriori

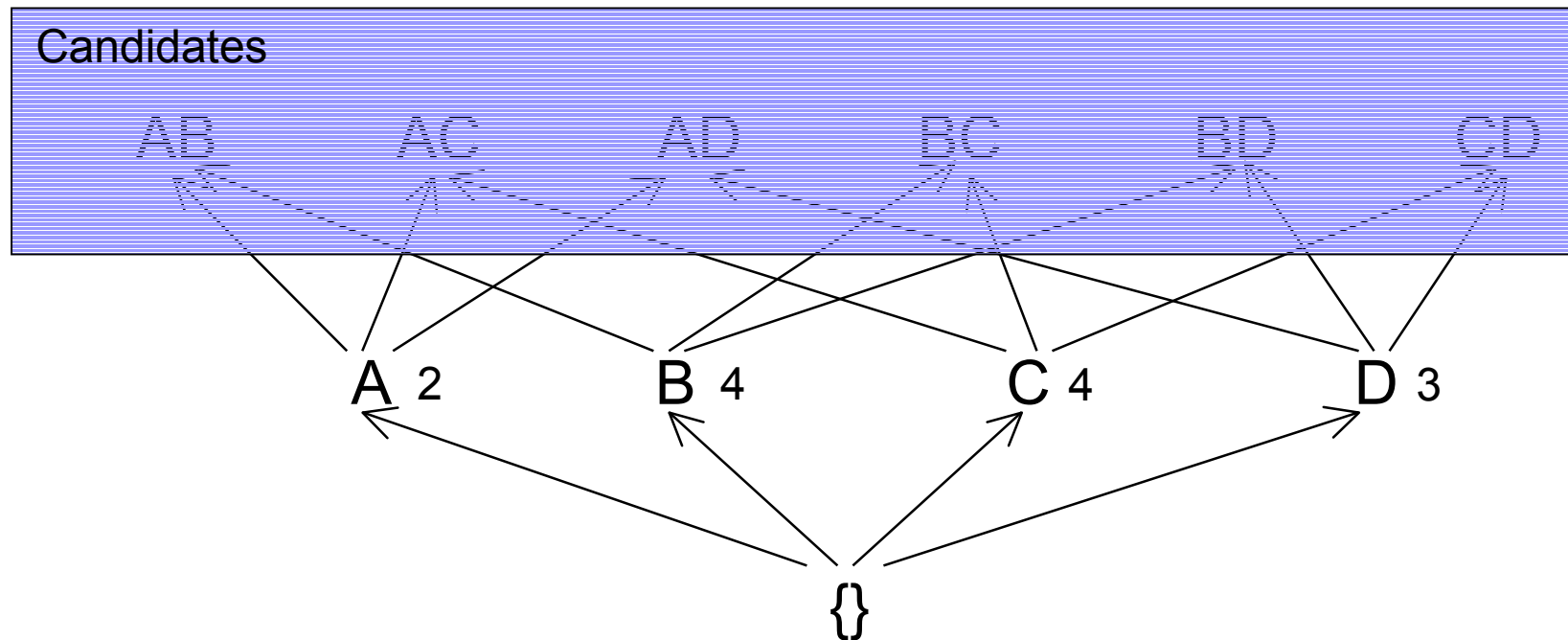
minsup=2



TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

Apriori

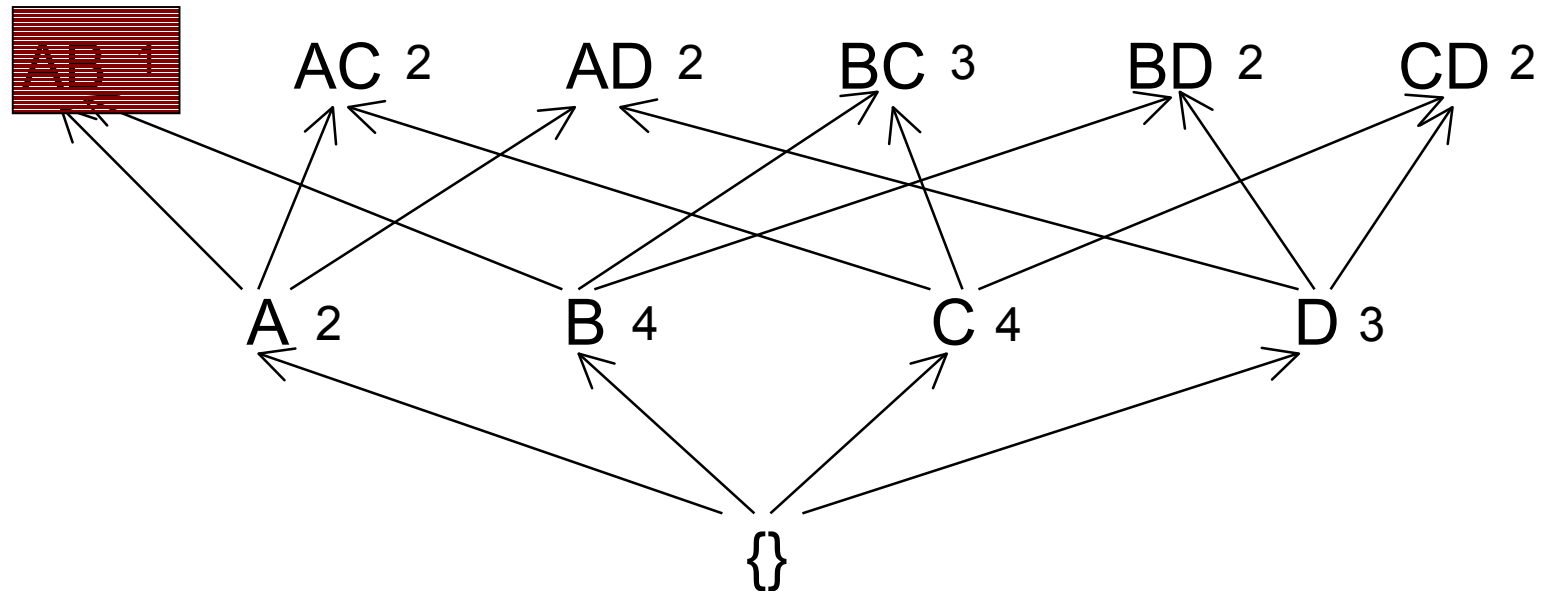
minsup=2



TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

Apriori

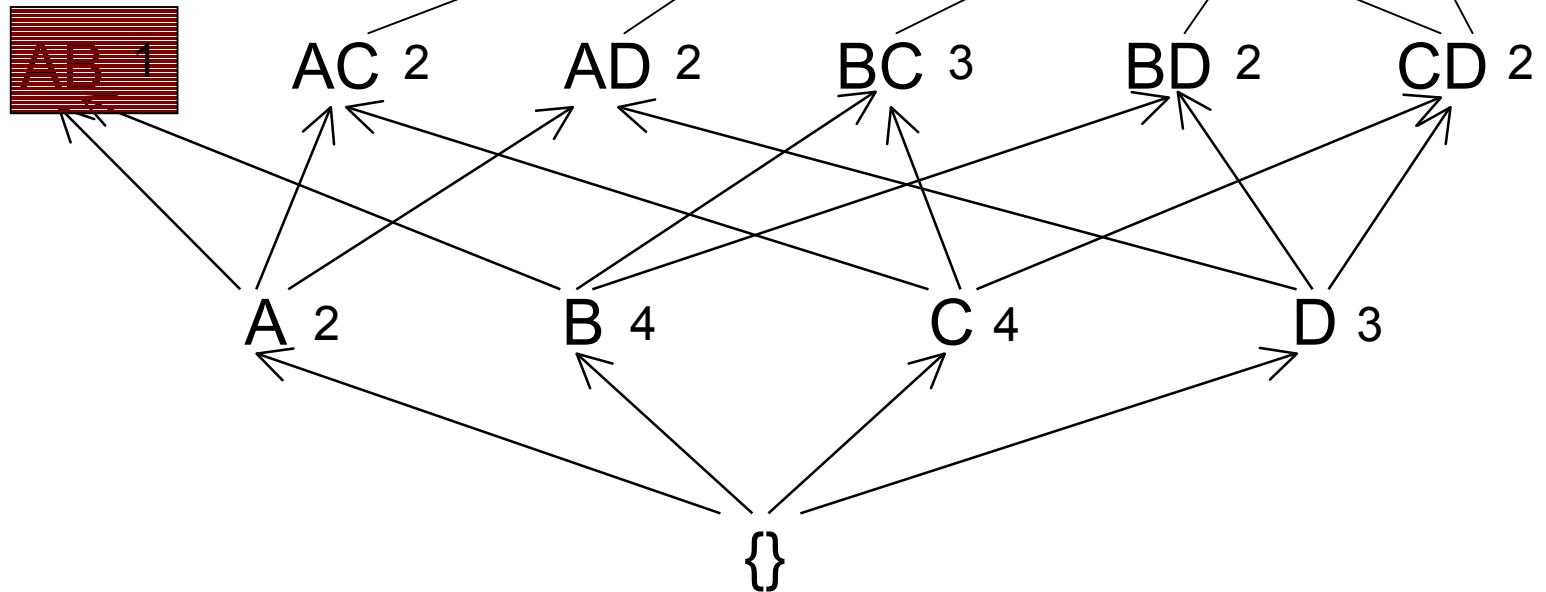
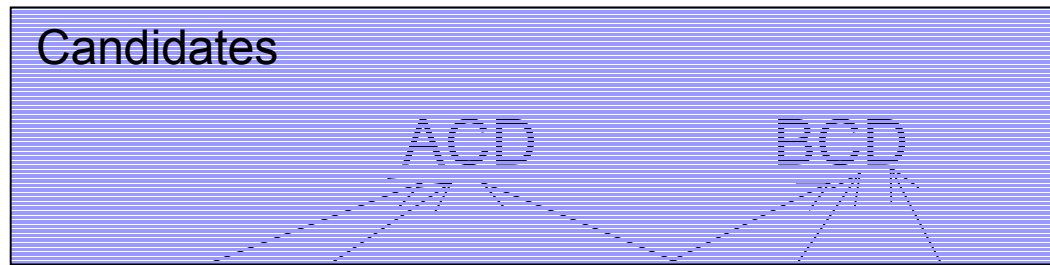
minsup=2



TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

Apriori

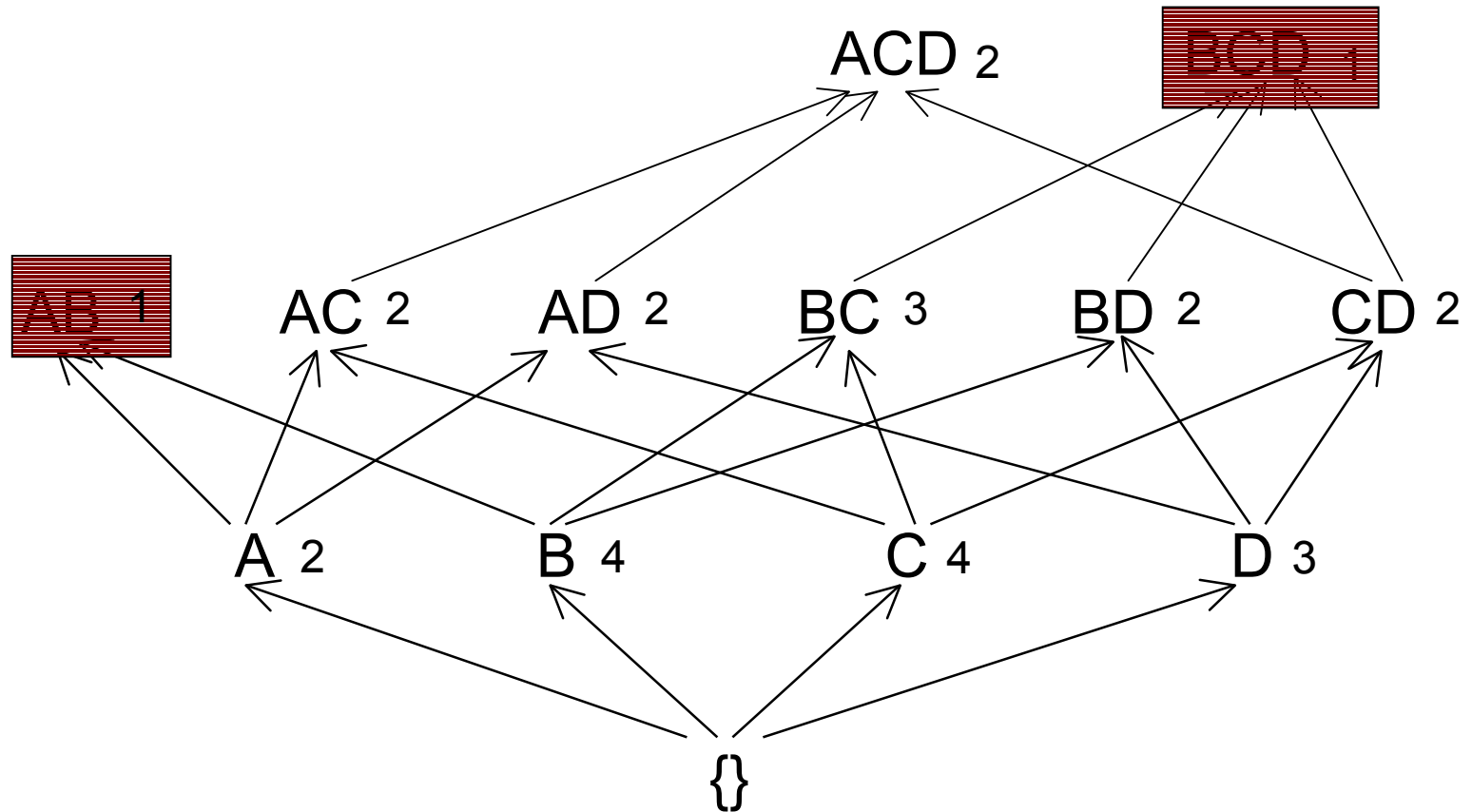
minsup=2



TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

Apriori

minsup=2



Depth-First Algorithms

Find all frequent itemsets

TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

Find all frequent itemsets, with D

TID	A	B	C	
3	1	0	1	
4	1	1	1	
5	0	1	0	

Find all frequent itemsets, without D

TID	A	B	C	
1	0	1	1	
2	0	1	1	
3	1	0	1	
4	1	1	1	
5	0	1	0	

Depth-First Algorithm

A,B,C,D
 AD, BD, CD
 ACD
 AC, BC

TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

DB[D]

TID	A	B	C
3	1	0	1
4	1	1	1
5	0	1	0

DB[CD]

TID	A
3	1
4	1

DB[BD]

TID

DB[C]

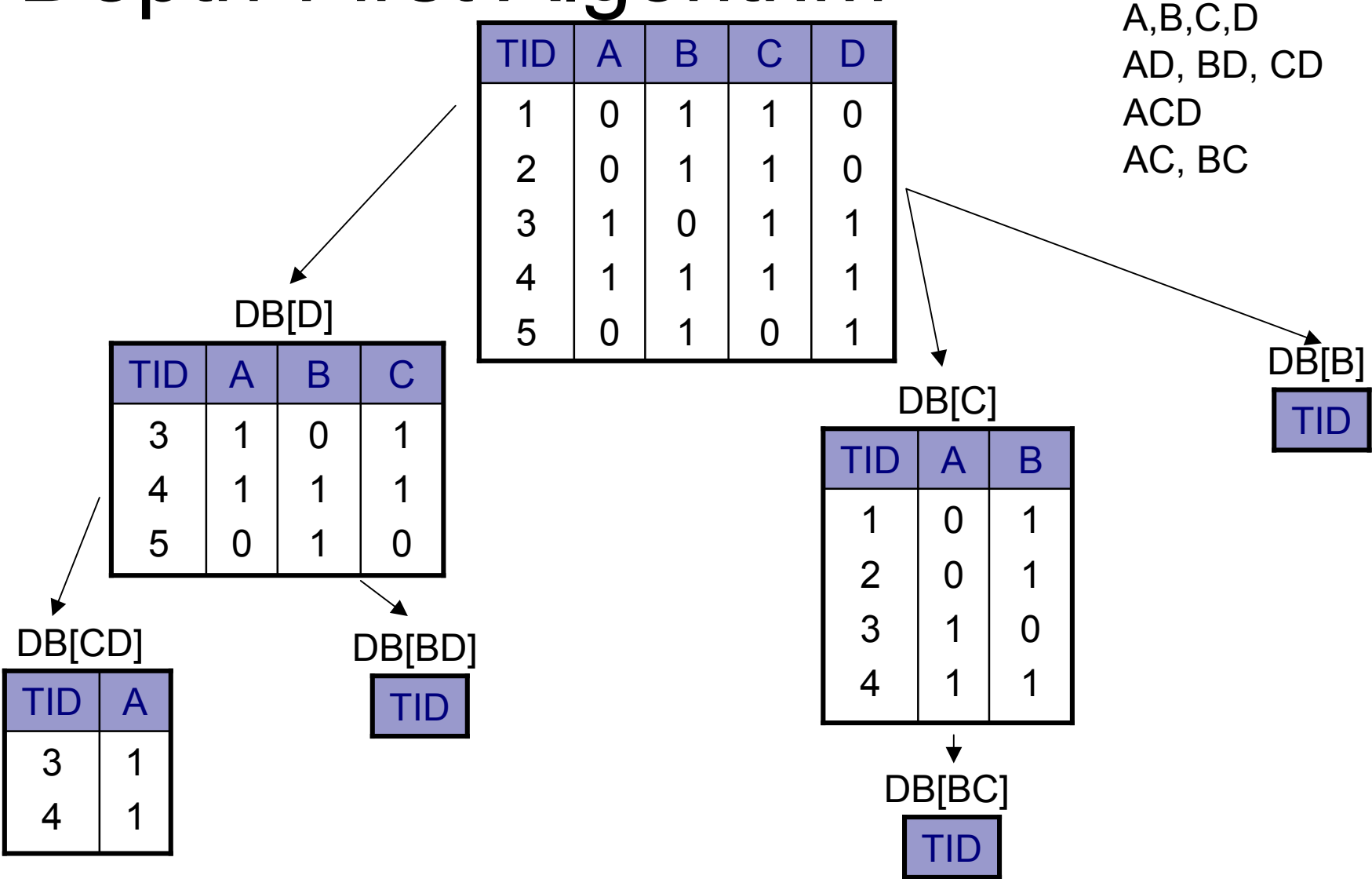
TID	A	B
1	0	1
2	0	1
3	1	0
4	1	1

DB[BC]

TID

DB[B]

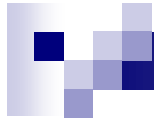
TID





Breadth-First vs Depth-First

- Depth-first outperforms breadth-first
 - Number of frequent itemsets is very high
 - Database is relatively small
- Breadth-first outperforms depth-first
 - Number of frequent sets is small
 - Database is large
- Differences usually very small



Summary

- Frequent Itemset Mining
- Algorithms
- Constraint Based Mining
- Condensed Representations



Mining With Constraints

- Reduce output size, user sets focus
 - itemsets of size > 5
 - sets of products with cost less than 10 EUR
 - sets that contain A, B, or C.
 - sets that are frequent in dataset D_1 , but infrequent in D_2



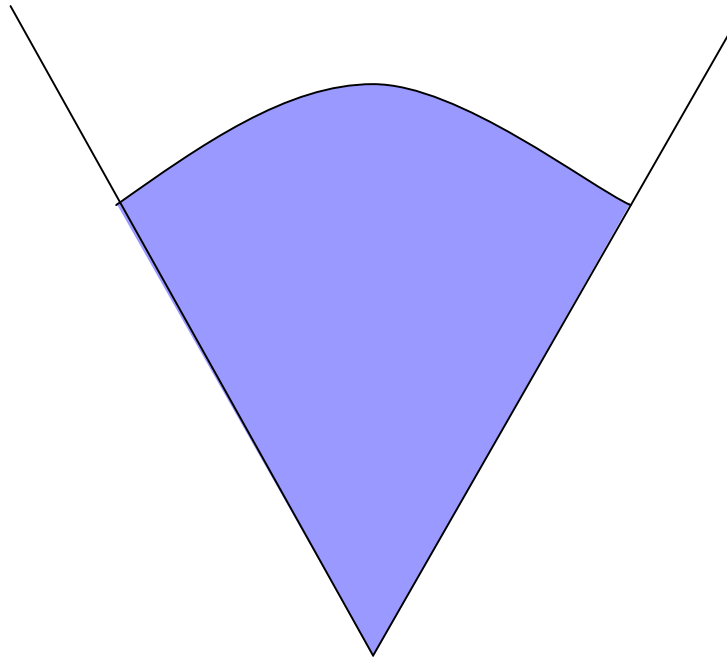
Mining With Constraints

- Types of constraints
 - (Anti-)Monotone,
 - Succinct
 - Convertible
- Two Approaches
 - Pushing constraints into the mining algorithm
 - Changing the Database



Types of Constraints

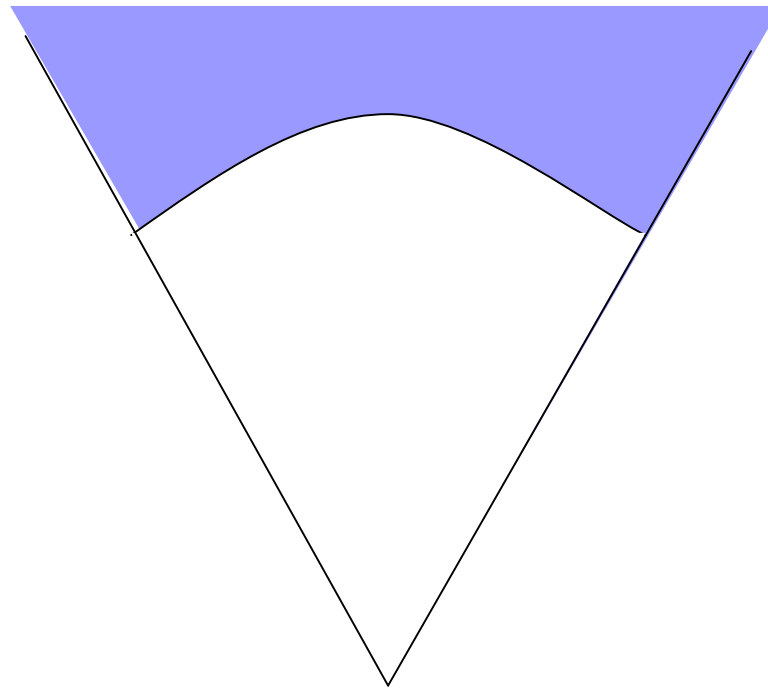
- Anti-monotone
 - Support, size < 10 , ...



Types of Constraints

- Monotone

- Cost > 10EUR, Contains A, B, or C, ...





Types of Constraints

■ Succinct

□ Can be expressed using minus and union on a fixed number of powersets

■ E.g., Contains A or B, but not C: $2^{I-C} - 2^{I-AB}$

□ Can be generated efficiently

■ Convertible anti-monotone

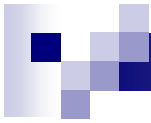
□ Anti-Monotone w.r.t. prefix-order

■ E.g. $\text{avg}(I.\text{price}) < 10$ EUR when ordered ascending by price.

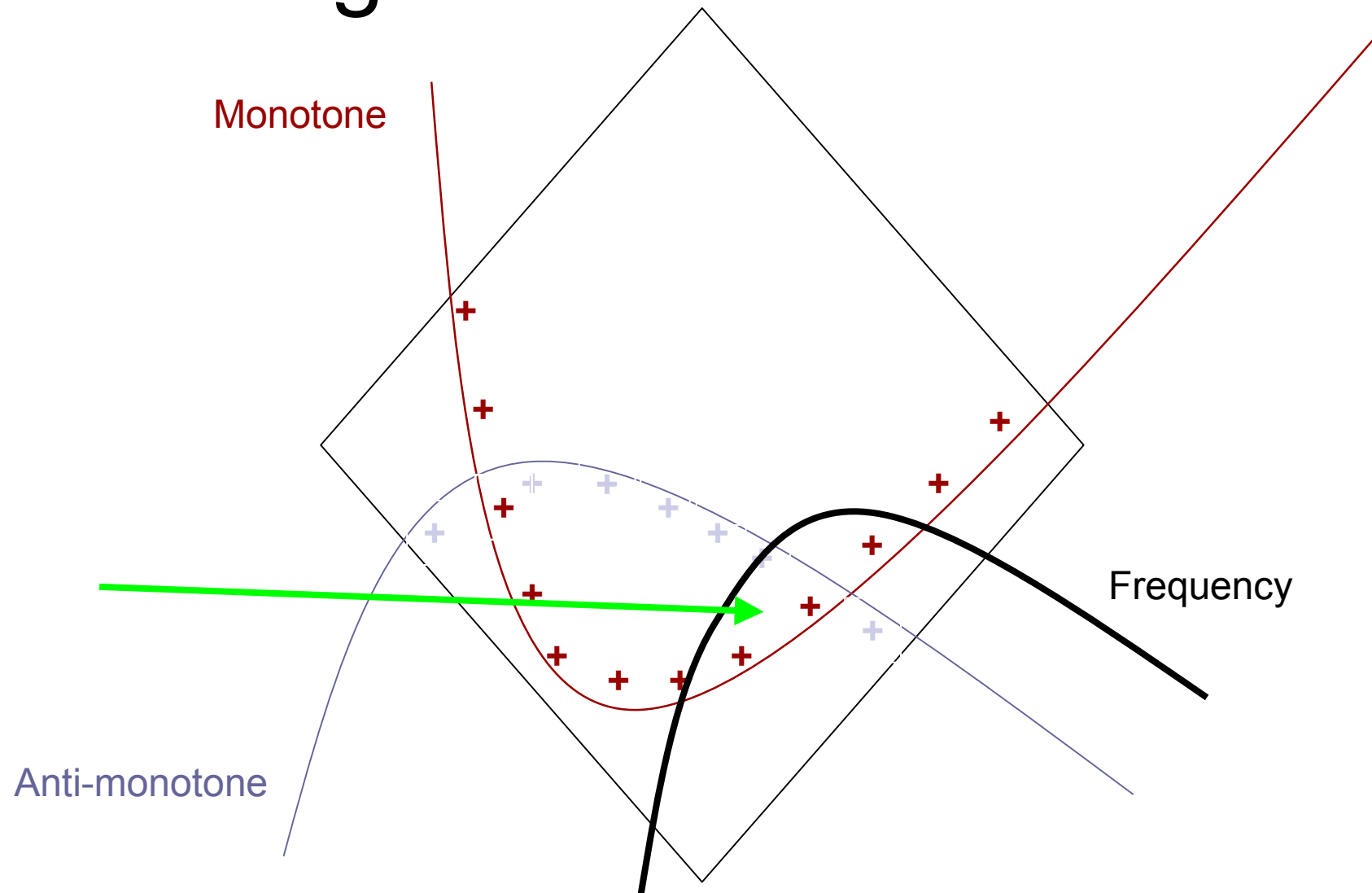


Mining With Constraints

- Two approaches:
 - Pushing constraints deep in data mining algorithm
 - Changing database such that
 - Support of itemsets satisfying the constraint does not change
 - The support of itemsets that do not satisfy the constraint decreases



Pushing Constraints





Pushing Constraints

- Trade-off
 - Pushing monotone constraints
 - vs. anti-monotone pruning
- Not always better to push monotone constraints
 - E.g. Size > 10 ...



Changing the Database

■ ExAnte Algorithm

- Exploit Monotone and Anti-monotone constraints
- A transaction that does not satisfy a monotone constraint will not contribute to any itemset satisfying the constraints
 - E.g. constraint “size > 10”: every transaction of size < 10 can be thrown away!

Changing the Database

minsup = 3
size ≥ 4

anti-mon.
monotone

ID	A	B	C	D	E	F	G	H	I
1	1	1	0	0	0	0	1	0	0
2	1	0	1	0	0	1	0	1	0
3	1	1	1	1	0	0	1	0	1
4	1	0	0	0	1	1	0	1	1
5	1	1	1	1	1	0	0	0	1
6	1	0	1	1	0	0	0	0	1
7	0	1	0	1	0	0	1	0	0



Summary

- Frequent Itemset Mining
- Algorithms
- Constraint Based Mining
- Condensed Representations



Condensed Representations

- Sometimes, the output of frequent set mining remains too large:
 - Huge number of items
 - Highly correlated
 - High support items
- Hence, instead of mining all itemsets
 - Condensed representation



Condensed Representations

- Closed sets

- Divide frequent itemsets into equivalence classes
- Two itemsets are equivalent if they occur in the same transactions
- Closed set: maximal element in an equivalence class

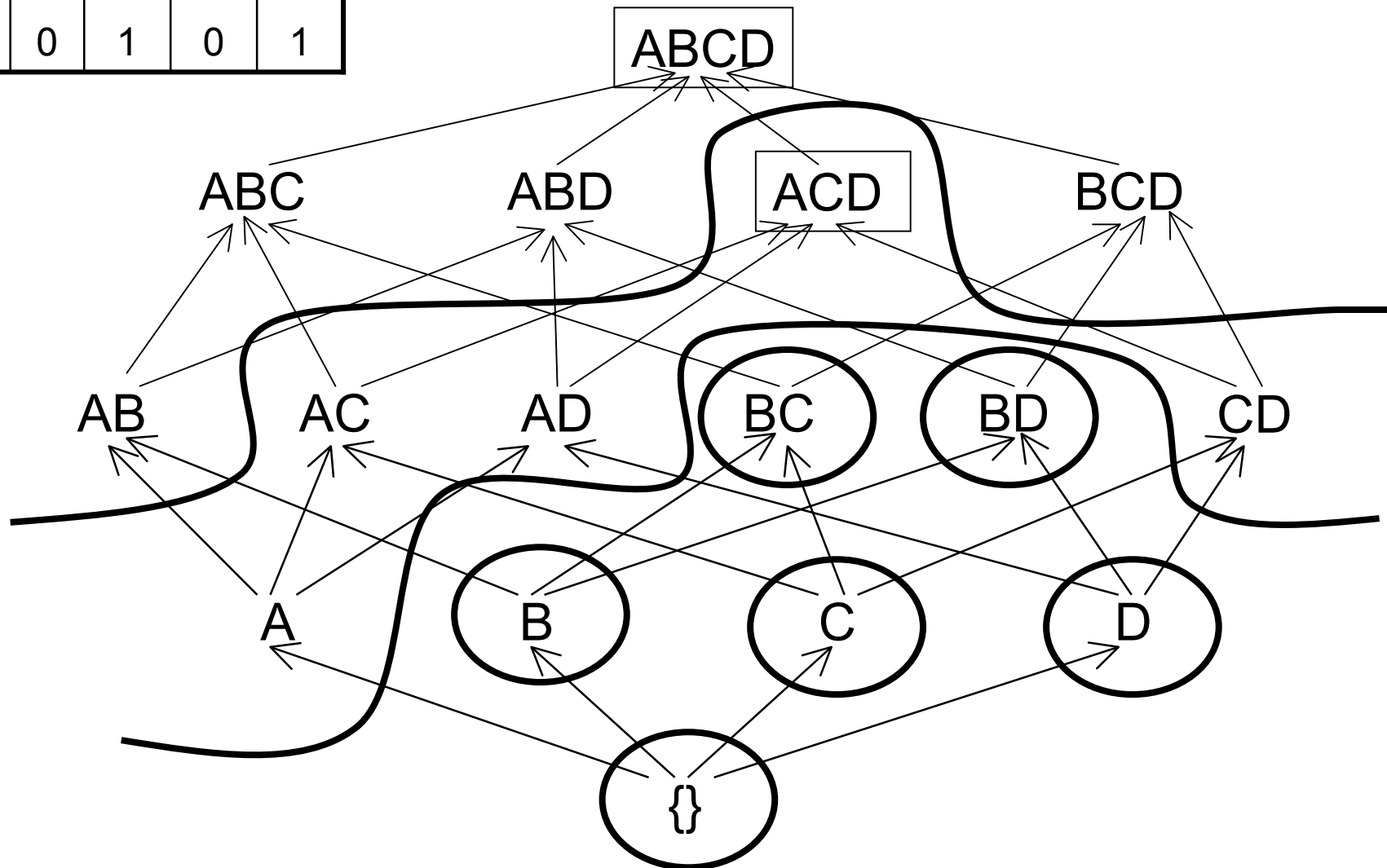


Closed Itemsets

- All sets in the same equivalence class have the same support
 - Occur in the same transactions
- Maximal element in an equivalence class is unique
 - If two itemsets occur in the same transactions, then so does their union

TID	A	B	C	D
1	0	1	1	0
2	0	1	1	0
3	1	0	1	1
4	1	1	1	1
5	0	1	0	1

Closed Itemsets





Closed Itemsets

- Has nice mathematical properties
 - Closed sets form a lattice
 - Galois connection
- Efficient algorithms to find them
- Based on the closed sets, it is easy to find the support of the other itemsets.



Closed Itemsets

- Interesting class of patterns
 - Maximal frequent itemsets are closed sets
 - Highest correlation between items
 - Strongest association rules
- Significant reduction of number of itemsets
 - Especially with small number of large transactions



Non-Derivable Itemsets

- Based on redundancies
 - How do supports interact?
- What information about unknown supports can we derive from known supports?
 - Concise representation: only store relevant part of the supports



Redundancies

- Agrawal et al. (Monotonicity)
 - $\text{Supp}(AX) \leq \text{Supp}(A)$
- Boulicaut et al., Lakhal et al. (Free sets)
 - **If** $\text{Supp}(A) = \text{Supp}(AB)$ (Closed sets)
Then $\text{Supp}(AX) = \text{Supp}(AXB)$

Redundancies

- Bayardo

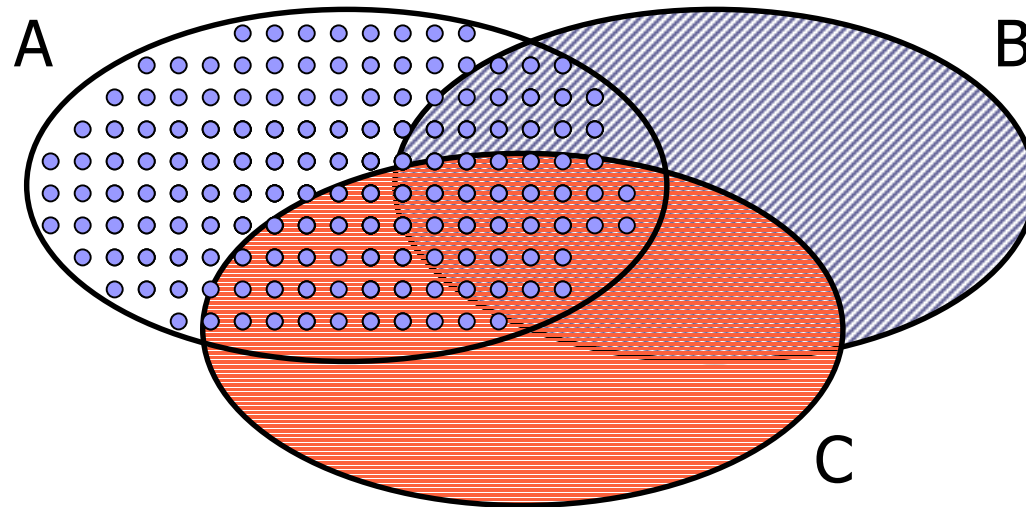
(MAXMINER)

- $\text{Supp}(ABX) \geq \text{Supp}(AX) - \frac{\text{Supp}(X) - \text{Supp}(BX)}{\text{drop}(X, B)}$

- Bykowski, Rigotti (Disjunction-free sets)

if $\text{Supp}(ABC) = \text{Supp}(AB) + \text{Supp}(AC) - \text{Supp}(A)$, then $\text{Supp}(ABCX)$ can be derived from ABX, ACX, AX

The Inclusion – Exclusion Principle



$$\begin{aligned} |A \cup B \cup C| &= |A| + |B| + |C| \\ &\quad - |A \cap B| - |A \cap C| - |B \cap C| \\ &\quad + |A \cap B \cap C| \end{aligned}$$



Deduction Rules via Inclusion-Exclusion

- Let A, B, C, \dots be items
- Let A' correspond with the set
 $\{ \text{transaction } t \mid t \text{ contains } A \}$
- $AB' = A' \cap B'$
- Then: $\text{Supp}(ABC) = |ABC'|$



Deduction Rules via Inclusion-Exclusion

- Inclusion-exclusion principle:

$$\begin{aligned} |A' \cup B' \cup C'| &= |A'| + |B'| + |C'| \\ &\quad - |AB'| - |AC'| - |BC'| \\ &\quad + |ABC'| \end{aligned}$$

Thus, since $|A' \cup B' \cup C'| \leq n$,

$$\begin{aligned} \text{Supp}(ABC) &\leq s(AB) + s(AC) + s(BC) \\ &\quad - s(A) - s(B) - s(C) + n \end{aligned}$$

Complete Set for Supp(ABC)

$$0 \quad s_{ABC} \geq 0$$

$$s_{ABC} \leq s_{AB}$$

$$1 \quad s_{ABC} \leq s_{AC}$$

$$s_{ABC} \leq s_{BC}$$

Monotonicity

Free, Closed

$$s_{ABC} \geq s_{AB} + s_{AC} - s_A$$

$$2 \quad s_{ABC} \geq s_{AB} + s_{BC} - s_B$$

$$s_{ABC} \geq s_{AC} + s_{BC} - s_C$$

Disjunction-Free

$$3 \quad s_{ABC} \leq s_{AB} + s_{AC} + s_{BC} - s_A - s_B - s_C + n$$



Derivable Itemsets

Given: $\text{Supp}(I)$ for all $I \subset J$

Lower bound on $\text{Supp}(J) = l$

Upper bound on $\text{Supp}(J) = u$

- Without counting : $\text{Supp}(J) \in [l, u]$
- J is a **derivable itemset** (DI) iff $l = u$
 - We **know** $\text{Supp}(J)$ **exactly** without counting!



Derivable Itemsets

J derivable itemset:

- No need to **count** Supp(J)
- No need to **store** Supp(J)
 - We can use the deduction rules
- Concise representation:
$$C = \{ (J, \text{Supp}(J)) \mid J \text{ not derivable from } \text{Supp}(I), I \subset J \}$$



Derivable Itemsets

Theorem (Monotonicity)

If $J \subset K$, J derivable, then K derivable.


Moreover:

The width of the interval for $J \cup \{A\}$ is at most half the size of the interval for J .



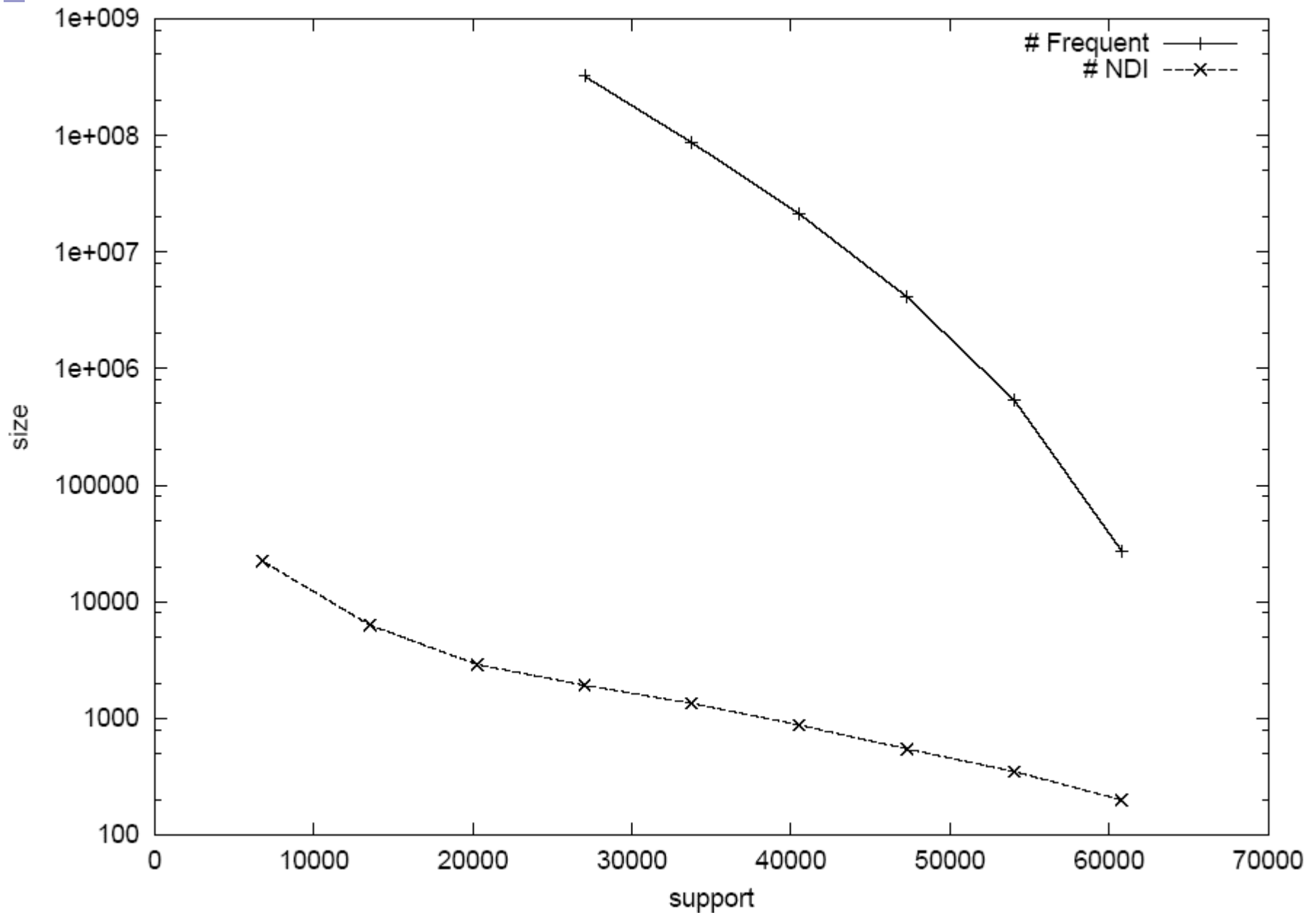
IV. Evaluation --- Theoretical

- Interval widths decrease exponentially
 - Half each step
- Non-derivable itemset can never be larger than $\log(|\text{Database}|)$
 - Independent of sparse, dense, ...

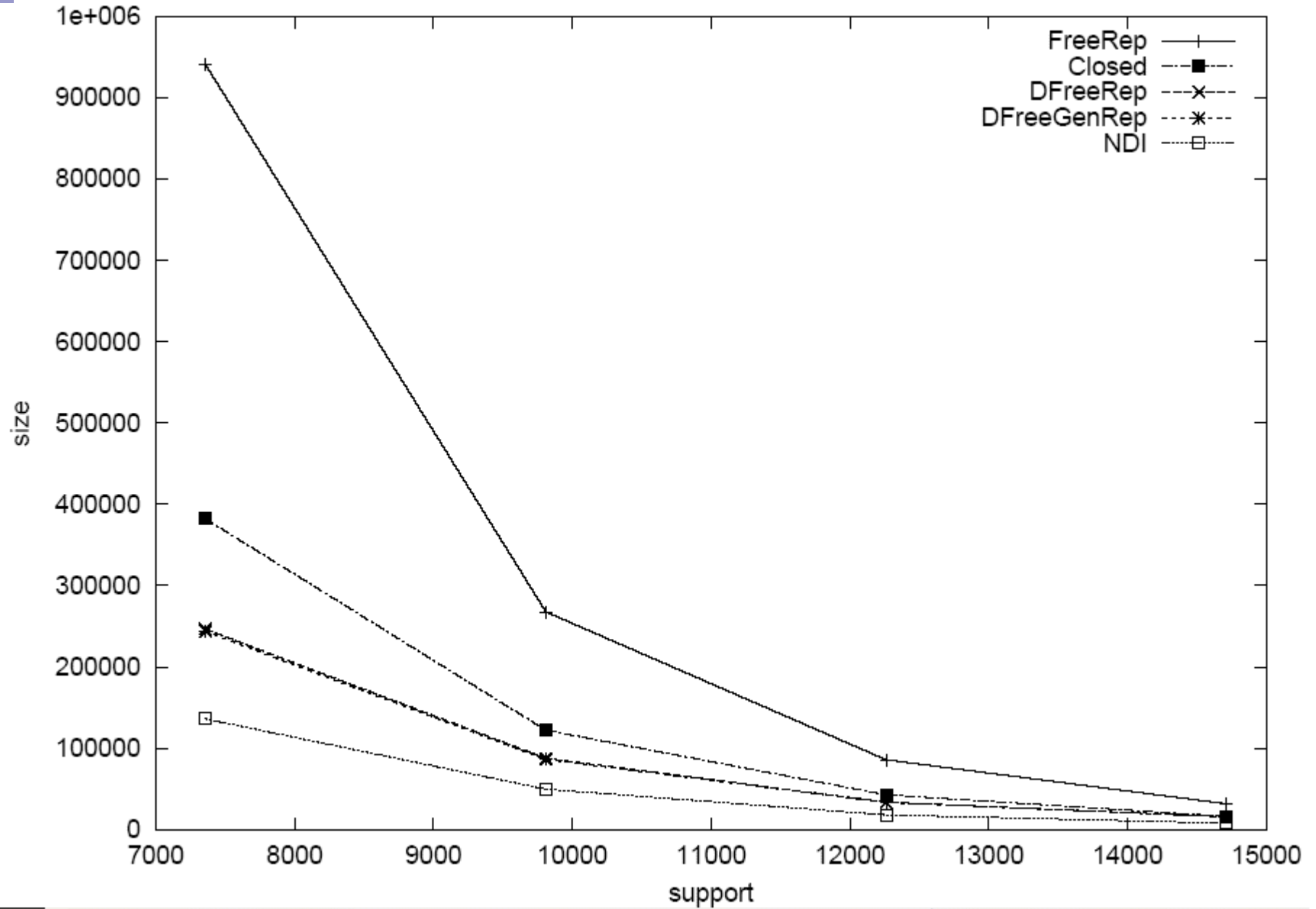


Evaluation --- Empirically

- Size NDI vs. frequent itemsets
- Comparison with Other Concise Reps



PUMSB



PUMSB



Evaluation

- Number of frequent NDIs considerable smaller than number of frequent itemsets
- Algorithm is efficient
 - Calculating NDI + deducing DIs often outperforms Apriori



Condensed Representations

- Many other representations
 - Free sets
 - Disjunction-free sets
 - Generalized disjunction-free sets
 - ...
- Closed sets and NDIs provable the smallest ones



Conclusion

- Depth-first vs Breadth-first algorithms for FIM
- Constraint mining to incorporate user focus
 - Pushing constraints vs changing database
- Condensed Representations
 - Closed sets
 - Non-Derivable Itemsets



Topics Not Covered ...

Parallel algorithms for FIM

Incremental FIM

Generalized, Quantitative, Multi-level, Fuzzy ARs

Coupling FIM with RDBMS

Privacy Preserving ARM

Computational Complexity Results

Inverse mining problem

Emerging Patterns, jumping emerging patterns

Dependency value, X^2

Lift, gain

Block support, tilings,

...